

S P E C I F I C A T I O N

SELF-DETERMINING COMMAND PATH ARCHITECTURE

BACKGROUND OF THE INVENTION

Field of the Invention

5 The invention relates to software design and particularly to a mechanism for software design for real-time embedded systems that facilitates deterministic behavior by providing data objects that include path objects that permit processing nodes to be defined
10 independently of special handling requirements for each data object.

Background

15 Object-Oriented (O-O) design is a method of design that 1) results in the decomposition of a problem in terms of objects and 2) employs separate models for the static (object structures) and dynamic (process architecture) design of a system. In the context of the object model, an "object" is a programmatic entity that
20 possesses state, behavior, and identity. The term "class" is applied to a definition of the common structure and behavior of a set of objects of a given type.

The object model has four major elements;
abstraction, encapsulation, modularity, and hierarchy.

Abstraction allows the separation of an object's essential
behavior from its implementation. Encapsulation allows the
5 details of an object that do not contribute to its
essential characteristics to be hidden from view.

Modularity allows the clustering of logically-related
abstractions, thereby promoting reuse. Hierarchy allows
the commonality of objects to be leveraged by providing for
10 the inheritance of the behavior of another object or
including other objects to achieve a desired behavior.

Each object within an O-O system is defined by an
interface and an implementation. A software client
external to an object depends completely on its interface
15 and not the details of its implementation. The
implementation of an object provides the mechanisms and the
details that define its behavior. O-O programs are
collections of objects that relate to each other through
their interfaces.

20 In a sense, each object is a "black box." Its
interface consists of messages that the black box sends and
receives. Objects actually contain code (sequences of
computer instructions) and data (information which the

instructions operate on). Traditionally, code and data have been kept apart. For example, in the C language, units of code are called "functions," while units of data are called "structures." Functions and structures are not formally connected in C. A C function can operate on more than one type of structure, and more than one function can operate on the same structure. This is not true for O-O software. In O-O programming, code and data are merged into a single indivisible thing -- an object. A programmer using an object should not need to look at the internals of the object once the object has been defined. All connections with the object's internal programming are accomplished via messages; i.e., the object's interface.

A computer typically runs several processes, or "threads," simultaneously. Each thread has its own stream of computer instructions and operates independently of other threads. The stream of computer instruction would typically be encapsulated in one or more objects, termed "processing objects." The operating system is responsible for allocating CPU (central processing unit) time to each thread based on priorities and other criteria.

In defining a process architecture some O-O programs use an approach, called "pipelining," that uses

queues to mediate the transfer of a data object among
processing objects. Several process threads may be created
and each is associated with an input queue. In a pipelined
system, a thread gets data objects from its input queue,
5 processes the data object, and then places it in another
queue (some other thread's input queue).

What happens to a data object during the course
of processing by a processing object may determine what
queue the data object has to be sent to after being
10 processed by the processing object. That means, the
processing object has to have the logic built into it to
know which processing object's queue to place the data
object into after it is finished processing the data
object.

15 O-O design and design patterns are also
applicable to real-time embedded systems, systems developed
for a real-world application such as a control mechanism
for a piece of equipment. Real-time systems like these are
required to satisfy precise timing deadlines, have
20 predictable response times, and must be stable when
overloaded. Such a system is said to be "deterministic" if
it can be guaranteed to respond by a certain time no matter
what happens. But high volume deterministic processing is

difficult due to the complex processing required for each
command (data object). Each command may require parsing,
interpretation, extrapolation, execution, status updating,
and feedback to the object that placed the data message or
5 data object in its queue. Trying to do all of this in one
step makes it difficult to meet hard timing deadlines.

Breaking the large task into smaller manageable tasks makes
it difficult to keep track of commands as they move through
the system. What is needed, is a design that provides an
10 architecture that allows the tasks to be cleanly separated
as in O-O systems and allows keeping track of data objects
as they move about the system.

15 SUMMARY OF THE INVENTION

The invention allows the creation of path objects
to connect processing objects logically in a system. Any
number of path objects can be created. The design isolates
the command objects and the processing objects from any
20 knowledge of the command object's paths through the system.
The path object itself is an organization of the available
queues in the system. At each step, there is a queue for

the normal path, one for the error path, and any number of other queues for other processing outcomes.

Each processing object in a path is associated with an input queue. Each command object is associated with a path object. When a command is retrieved from the processing object's input queue, it is processed and then sent to the next queue in its path (for example, either normal path or error path). This allows each processing object to focus on its task without being encumbered by details of the command or its task in the system. This is particularly useful in addressing high-volume deterministic commands because the schedule and execution time of the critical processing objects can be tightly controlled. Other steps in the command path can be processed at a lower priority.

In essence, in the environment of an O-O system in which data is transferred between threads by queues, the invention provides that the destination queue for a given data object depends on a path object defined for the data object. In an embodiment, each instance of a data object corresponds to a path object. When a processing object is finished processing the data object, it consults the path object to determine which queue to put the data object in.

Thus, if any change in the command path needs to be made, it can be defined in the path object, rather than by making changes in the processing objects that handle the data objects.

5 The invention will be described in connection with certain preferred embodiments, with reference to the following illustrative figures so that it may be more fully understood. With reference to the figures, it is stressed that the particulars shown are by way of example and for
10 purposes of illustrative discussion of the preferred embodiments of the present invention only, and are presented in the cause of providing what is believed to be the most useful and readily understood description of the principles and conceptual aspects of the invention. In
15 this regard, no attempt is made to show structural details of the invention in more detail than is necessary for a fundamental understanding of the invention, the description taken with the drawings making apparent to those skilled in the art how the several forms of the invention may be
20 embodied in practice.

BRIEF DESCRIPTION OF THE DRAWING

Fig. 1 is a flow chart illustrating a process for a processing object to determine a destination queue responsively to a data object according to an embodiment of the present invention.

Fig. 2 is a block diagram illustrating information flow between objects according to an embodiment of the invention.

Fig. 3 is a UML (unified modeling language) diagram illustrating the classes of an embodiment of the present invention.

Fig. 4 is a UML sequence diagram illustrating one example execution path for an embodiment of the invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

Referring to Fig. 1, processing objects in a program accept data from a queue in a first step S100. The processing object performs a data operation on the data object S110. The processing object determines a status relating to the processing of the data object, for example, whether the process was completed normally or with a fault in step S110. To determine a destination queue to which the data object is to be placed, the processing object

examines, or uses, data in the data object's path object in step S120. Then the processing object places the data object in the referenced queue in step S130.

The step of determining a status relating to the data object in step S110 is not essential to the invention. The status, according to an embodiment, may include normal and faulty results of processing.

Referring to Fig. 2, in an embodiment of the invention, an Nth data object 110 is placed in a queue 141 for a processing object 1 100. The processing object 1 100 processes the Nth data object 110 and queries an Nth data object's path object 115 which stores a table of queue indicators. In this embodiment, normal and faulty outcome states are defined for the processing object 1 100. When the processing object 1 100 determines the outcome of its internal processing, an indicator of this outcome points to an indicator 181, 182 stored on the Nth data object's path object 115 that refers to a particular queue to which the Nth data object 110 should next be placed. In the example, the status indicates normal and the queue 142 to which the selected indicator 181 points is that for processing object 2 120. The processing object 1 100 places the Nth data

object 110 in the queue 142 in response to the receipt of a message so indicating.

As a result of the above pipelined structure, the data object in this system determines its own destiny by its association with a path object. As a result, routing issues can be defined and adjusted in each path object rather than by programming this directly in the processing objects. Each processing object may be programmed to insure that a path object is queried to determine the destination of the data object once processing is complete. Note that a single path object can serve multiple data objects and vice versa according to the way in which the programmer chooses the package the information.

A condition, for example whether the processing object completed its process normally or anomalously, may not be present. The path object could simply carry a pointer respective of the current processing object indicating the destination queue. Also, the destination queue could be derived from a formula rather than simply static pointer. Further, the path data in the Nth data object's path object 115 can be incorporated directly in the data object itself.

Fig. 3 is a UML class diagram (static model) showing the structure of, and relationship between, objects in the system. Each rectangular box, e.g., path class 210, represents a class (potentially one or more objects). The box is divided horizontally into three sections, e.g., 211, 212, 213. A top section, e.g., 211, contains the class name. A middle section, e.g., 212, contains the class data; represented by the notation `[+ -][name]:[type]` where "+" and "-" indicates public and private data, respectively, "name" is an identifier for the unique instance of a given type of data, and "type" is the type of the data (possibly another class). The bottom section, e.g., 213, contains the class functions (computer code); represented by the notation `[+ -][func_name]([arg_N_name]:[type_N]):[func_type]` where "+" and "-" indicates public and private functions respectively, "func_name" is the function's name, "arg_N_name" identifies the Nth argument passed to the function, "type_N" is the type of the Nth argument, and func_type is the type returned by the function.

In Fig. 3 the Processor class 240 represents a processing object. It has an input queue and a function to process the command data. The Queue class 220 represents a

queue of commands. It has functions to return the next
command object in the queue and to add a command object to
the queue. The command class 230 represents the data that
the processing object instance of the Processor Class 240
5 will process. It keeps track of its position in the
command path using a step variable. It has one path and
accessor functions to return the next queue. The path
class 210 represents a unique path through the system. It
may be defined as an array of queue structures as detailed
10 in the note 200.

Fig. 4 is a UML sequence (dynamic model) showing
one example of the possible interactions between instances
of the objects defined in the class diagram. The boxes
arranged horizontally across the top of the figure are
15 objects (instances of the same classes from Fig. 3). The
notation is [name]:[type] where "name" uniquely identifies
and instance of "type", which is a class. The dashed line
descending from each object is its life line. The boxes
superimposed over the life lines represent activity of the
20 respective object. The horizontal lines between the
activity boxes represent messages between the objects, with
the arrow indicating the direction of flow. The text above

each message is the name of the function call pertaining to the message. Time flows from top to bottom in the diagram.

In Fig. 4 an unspecified event activates Node1 241. Subsequently Node1 241 makes a call (sends a message) to Que1 221, specifically GetNextCommand. Que1 221 returns the next available command object (data object), commandN 231 (a generic representation of any command in the system). Next, Node1 241 performs requisite processing on the command object. When processing is finished Node1 241 calls GetNextNormalQueue on commandN 231. CommandN 231 delegates the call to NextNormalStep on its path, pathN 251 (a generic representation of commandN's 231 path object), passing its step as an argument. PathN 251 returns the requested Queue object (Que2 222 in this example), in this case the next normal queue. Finally, Node1 241 calls AddToQueue on Que2 222, passing the current command object 231 as its argument. This ends the sequence of events for the diagram in Fig. 4. Typically the sequence would subsequently repeat.

Note that while in the embodiments described above, the path object is described as an object that is separate from the data object to which it relates, it is

clear that the path object may be incorporated within the data object.

It will be evident to those skilled in the art that the invention is not limited to the details of the foregoing illustrative embodiments, and that the present invention may be embodied in other specific forms without departing from the spirit or essential attributes thereof. The present embodiments are therefore to be considered in all respects as illustrative and not restrictive, the scope of the invention being indicated by the appended claims rather than by the foregoing description, and all changes which come within the meaning and range of equivalency of the claims are therefore intended to be embraced therein.